

Meeting GECCOO Project

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Symbolic animation of JML specifications

Frédéric Dadeau

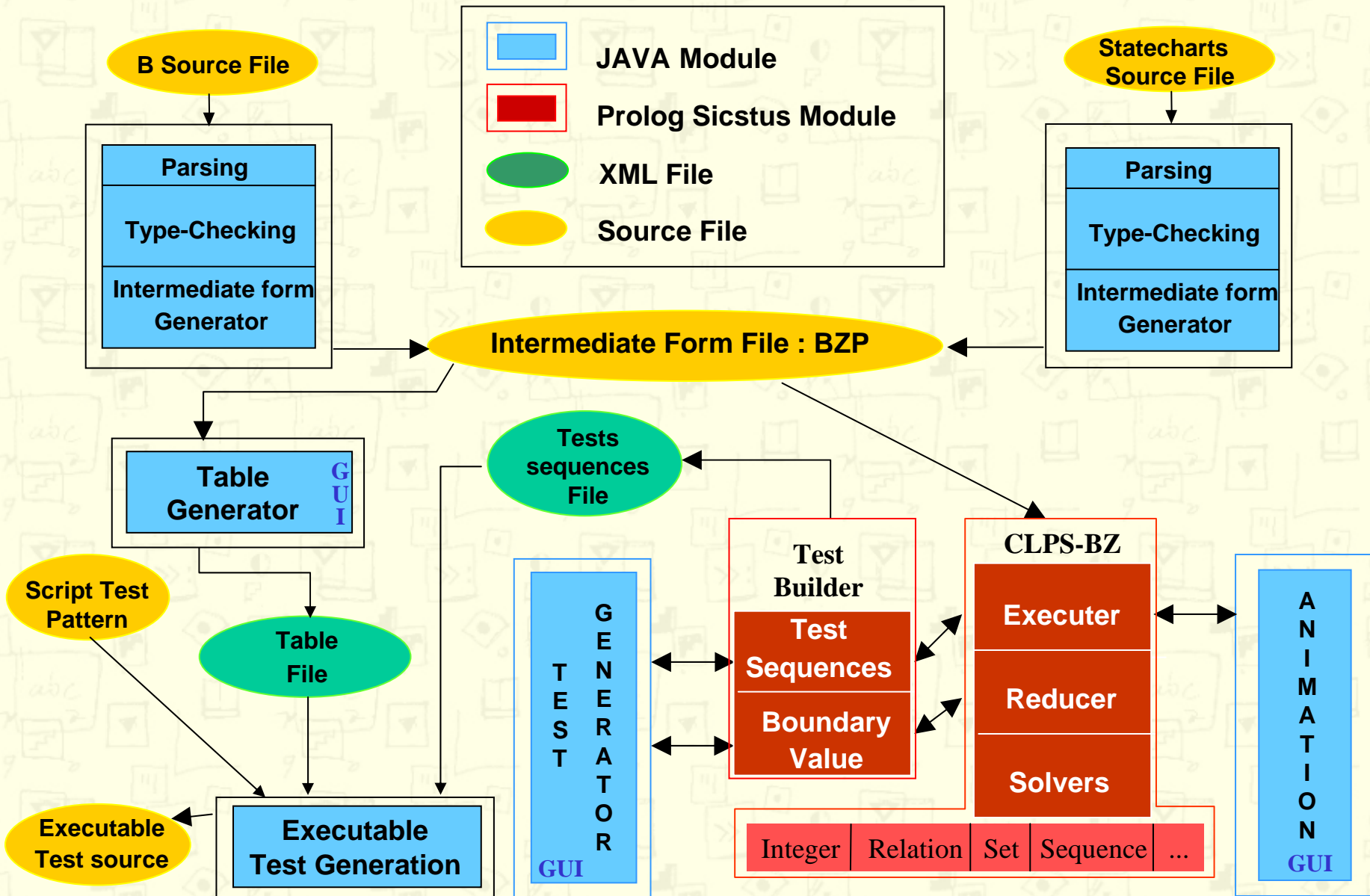




Introduction

- ① BZ-Testing-Tools: symbolic animator and model-based test generator based on constraint solving technologies
- ① Work: symbolic animation of JML specification, based on the model described in the JML annotations
- ① Motivation: use constrained animation to simulate the execution of the formal model and ensure its validity

Architecture





Outline

- ① The BZP(E) format
- ① Symbolic animation
- ① Translation of JML specifications to BZPE
- ① Checking properties
- ① The animator
- ① Future work



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The BZP(E) format

- ① Intermediate format of the BZ-Testing-Tools environment
- ① Pre / post-condition format
- ① Animation realized through the PROLOG module « executer » based on the CLPS-BZ solver
- ① Uses finite data structures



The BZP(E) format

🌀 PROLOG syntax:

specification(*spec-name*).

declaration(*spec-name*, **data-kind**, *data-name*, **data-type**).

operation(*spec-name*, *operation-name*).

predicat(*spec-name*, **pred-kind**, *pred-id*, *predicate*)

data-kind : static | variable | input(*op-name*) | output(*op-name*) | local(*op-name*)

data-type : atom | int | set(**data-type**) | pair(**data-type**, **data-type**)

pred-kind : static | invariant | initialisation | pre(*op-name*) | post(*op-name*)

The BZP(E) format

Semantics

- Use of a constraints store containing the set of variables values and their constraints
- Initialization
 - Add the « initialisation » predicate to the store
- Execution of an operation
 - Decomposition of the operation in behaviors (rewriting of disjunctions) ;
 - Parallel adding of different kind of predicates (precondition, postcondition, invariant) to the store.



The BZP(E) format

🌀 From BZP ... to BZPE

- BZPE is an extension of BZP
- Allows multiple specifications declaration
- References to data across specifications
- Allows translation of object concepts (ex. UML/OCL)



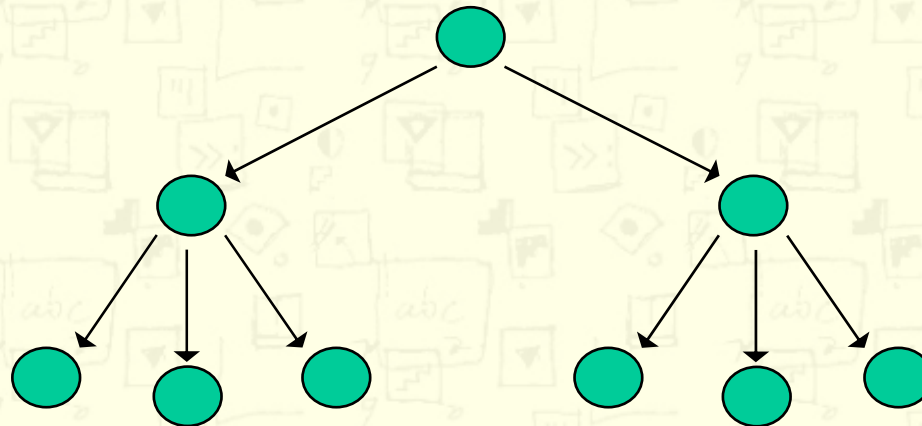
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Animation

« Classical » animation:

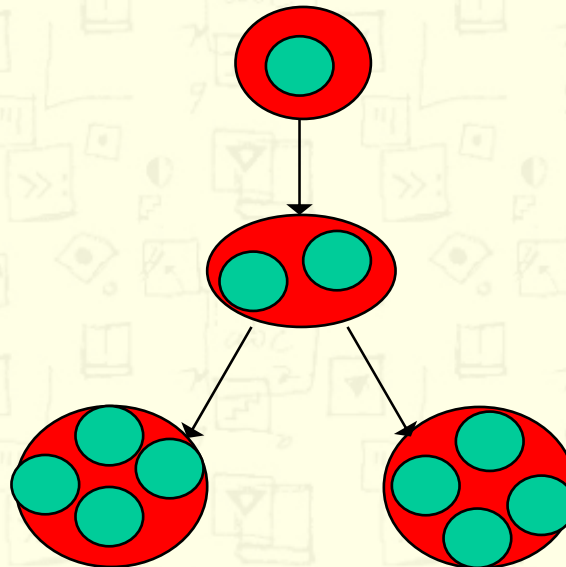
- Run through of a reachability graph
- Each node represents a valued state
- Each arc represents an activable behavior



Animation

🌀 Symbolic animation:

- Run through a constraint reachability graph
- Node are constraints stores
- A constraints store represents a finite set of valued states





Animation

🎯 What is a constraint store?

- Set of predicates giving informations on a variable's domain

- Each variable has a domain:

$$x \in [-2^{15}, 2^{15}]$$

- Each predicate added to the store reduces the variable domain:

$$x \in [-2^{15}, 2^{15}] \wedge x > 0 \rightarrow x \in [1, 2^{15}]$$

- If a predicates reduces a domain to the empty set, it becomes inconsistent:

$$x \in [1, 2^{15}] \wedge x = 0 \rightarrow x \in [] \rightarrow \text{inconsistency!}$$



Animation

🎯 How to create constrained states?

- With an unvalued initial state
- By leaving operation's input parameters unvalued



Animation

🌀 Advantages of symbolic animation:

- Dramatical decrease of the reachability graph size
- Propagation of the non-determinism of data
- The user is free to valuate input parameters or not



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Translation of JML specifications into BZPE

🌀 Translation of class structure

- Class → BZPE specification
 - With a set of possible instances (constant)
 - With a set of created instances (variable)
- Class attribute → BZPE variable
 - A total function mapping the set of created instances to the value
- Class method → BZPE operation
 - The class instance is an input parameter
 - Each method parameter → A BZPE operation parameter

Translation of JML specifications into BZPE

🌀 Translation of JML clauses

- Creation of new BZPE kind of predicates

- `jml_invariant` ;
- `jml_constraint` ;
- `jml_when (operation)`

interpreted by a specific « executer » ;

- Requires clause → BZPE operation pre-condition
- Ensures clause → BZPE operation post-condition

Translation of JML specifications into BZPE

```
public class date  
  int day ;  
  int month ;  
  int year ;  
  
  //@ invariant day > 0 &&  
  //@           day < 32 ;  
  
  /*@ public behavior  
  @ ensures  
  @   \result == \old(day) ;  
  @*/  
  
  public int getDay() {  
    return day ;  
  }
```

specification(date).

declaration(date, static, all_instances, set(atom)).
predicat(date, static, 1, all_instances = {d1,...,dN}).
declaration(date, variable, instances, set(atom)).
predicat(date, invariant, 2, instances \subseteq all_instances).
predicat(data, initialisation, 3, instances = {}).
predicat(date, invariant, 4, day \in instances \rightarrow IntRan)

predicat(date, jml_invariant, 5, $\forall x. x \in$ instances \Rightarrow day(x) > 0 & day(x) < 32).

operation(date, getDay).

declaration(date, input(getDay), this, atom).
declaration(date, output(getDay), result, int).
predicat(date, pre(getDay), 3, this \in instances).
predicat(date, post(getDay), 4, result = day(this)).



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Properties checking

⊙ How to check a **satisfiable** predicate:

- A constraint store represents a constrained state
- Addition of the predicate to the store:
 - If the store is still consistent: the predicate is satisfiable
 - If the store becomes inconsistent: the predicate is unsatisfiable
A valuation of the store presents a counter-example

⊙ Properties we want to be **satisfiables**:

- Pre-conditions (requires clause)
- Non divergence (diverges clause)
- Method activation after delaying (when clause)



Properties checking

🎯 How to check a **valid** predicate:

- A constraint store represents a constrained state
- Addition of the predicates negation to the store:
 - If the store is still consistent: the predicate is not valid
A valuation of the store presents a counter-example
 - If the store becomes inconsistent: the predicate is valid

🎯 Properties we want to be **valid**:

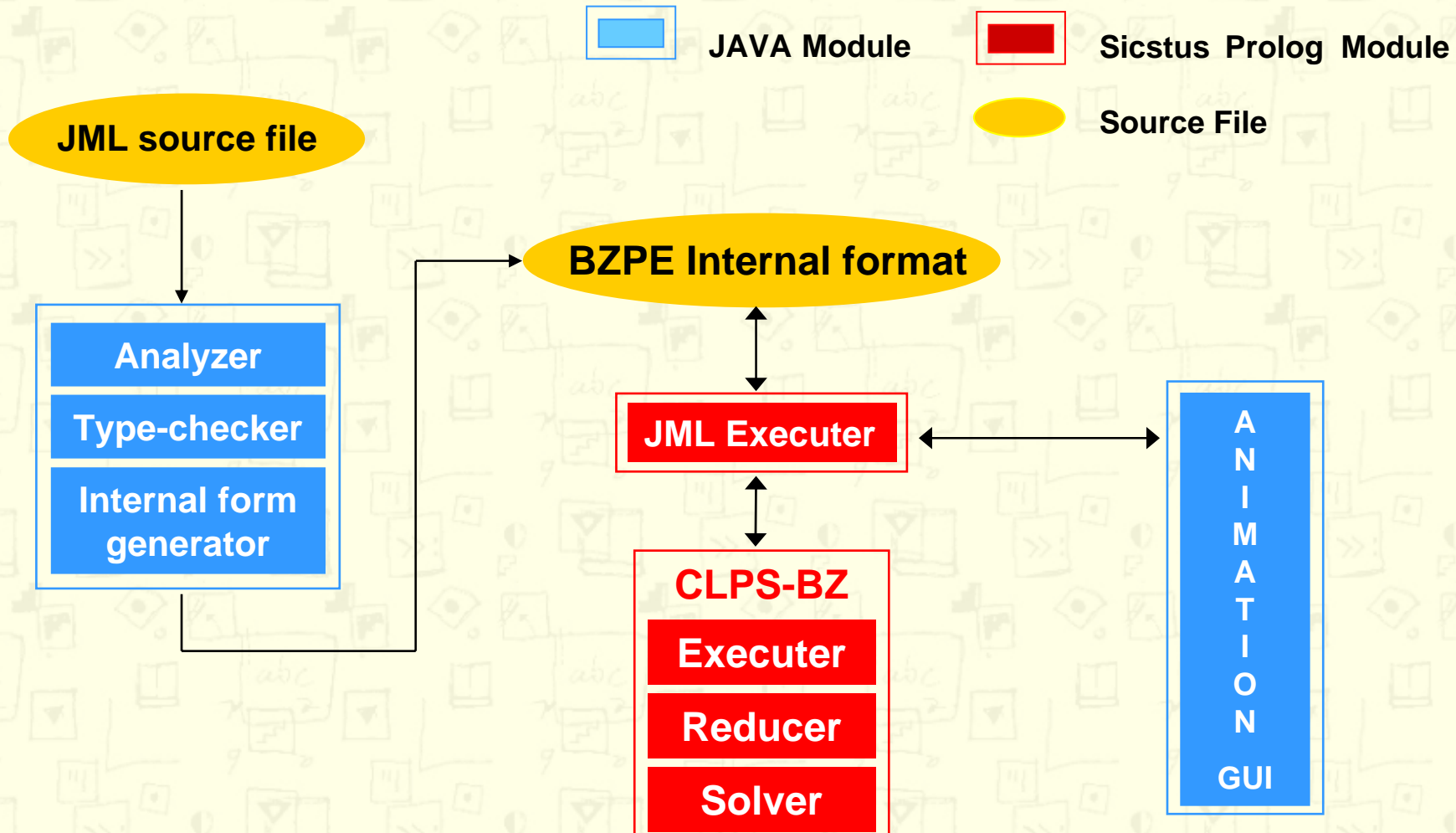
- Invariant
- History constraints



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The JML animator





The JML animator

- ⊙ Considers the model described by the JML annotations
- ⊙ Execution capabilities:
 - Valued or constrained animation ;
 - Activation of behaviors ;
 - Behaviors consistency ;
 - Method divergence avoiding ;
 - Method delaying ;
- ⊙ On-the-fly properties checking:
 - Class invariant;
 - History constraints;



The JML animator

Supported JML data types:

- integers (byte, short, int, char)
- arrays (multi-dimensionals)
- objects

Supported JML clauses:

- requires
- ensures
- diverges
- when
- invariant
- constraints
- initially



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Future work

🌀 Improving animator:

- Increasing JML support
 - Inheritance
 - Multiple specification files
 - Floats (waiting for solver)
- Translation
 - Method calls
- Type checking
- Improving GUI



Future work

🌀 From the constrained animation on:

- Manual / automatic simulation of the execution of a JML specification
- On-the-fly properties checking with reachables counter-examples production
- (Boundary) test cases generation
- Modelling errors detection
 - Too weak preconditions
 - Too strong / too weak invariant
 - ...